## f.root-servers.net

ISOC ccTLD Workshop Nairobi, Kenya, 2005

## The Basics

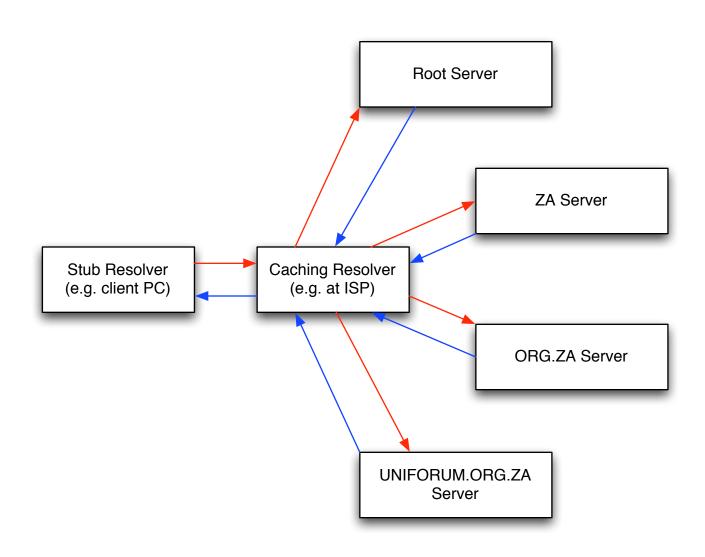
#### DNS

- The Domain Name System is a huge database of resource records
  - globally distributed, loosely coherent, scaleable, reliable, dynamic
  - maps names to various other objects
- The DNS allows people to use names to locate resources on the Internet, instead of numbers

## Components of the DNS

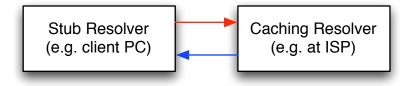
- A namespace
  - hierarchical, tree-like structure
  - labels separated by dots
- Nameservers
  - servers which respond to queries from clients, and make the data available
- Resolvers
  - clients which ask questions

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- Answers which are already in the cache can be returned directly, with no recursive lookup required
- Items expire from the cache when they become stale



#### Root Servers

- Every recursive nameserver needs to know how to reach a root server
- Root servers are the well-known entry points to the entire distributed DNS database
- There are 13 root server addresses, located in different places, operated by different people
- The root zone is published by IANA

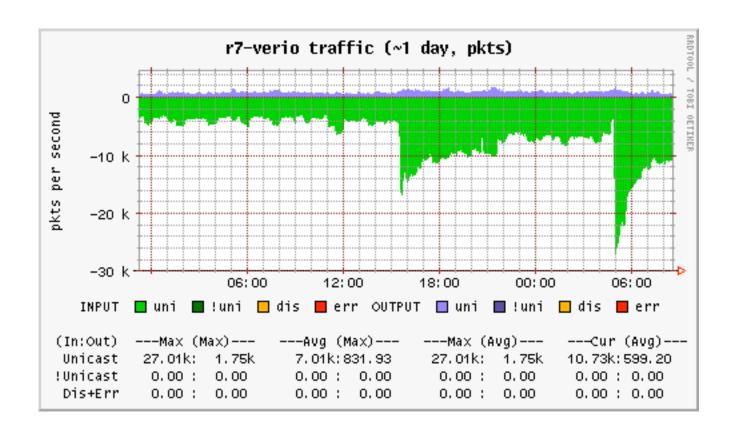
## The Root Servers

A.ROOT-SERVERS.NET	Verisign Global Registry Services	Herndon, VA, US
B.ROOT-SERVERS.NET	Information Sciences Institute	Marina del Rey, CA,
C.ROOT-SERVERS.NET	Cogent Communications	Herndon, VA, US
D.ROOT-SERVERS.NET	University of Maryland	College Park, MD, US
E.ROOT-SERVERS.NET	NASA Ames Research Centre	Mountain View, CA,
F.ROOT-SERVERS.NET	Internet Software Consortium	Various Places
G.ROOT-SERVERS.NET	US Department of Defence	Vienna, VA, US
H.ROOT-SERVERS.NET	US Army Research Lab	Aberdeen, MD, US
I.ROOT-SERVERS.NET	Autonomica	Stockholm, SE
J.ROOT-SERVERS.NET	Verisign Global Registry Services	Herndon, VA, US
K.ROOT-SERVERS.NET	RIPE	London, UK
L.ROOT-SERVERS.NET	IANA	Los Angeles, CA, US
M.ROOT-SERVERS.NET	WIDE Project	Tokyo, IP

## DNS Failure Modes

# Challenges on the Root

- There have been a number of attacks on the root servers
- Distributed denial of service attacks can generate a lot of traffic, and make the root servers unreachable for many people
- Prolonged downtime would lead to widespread failure of the DNS



# It's a Jungle Out There

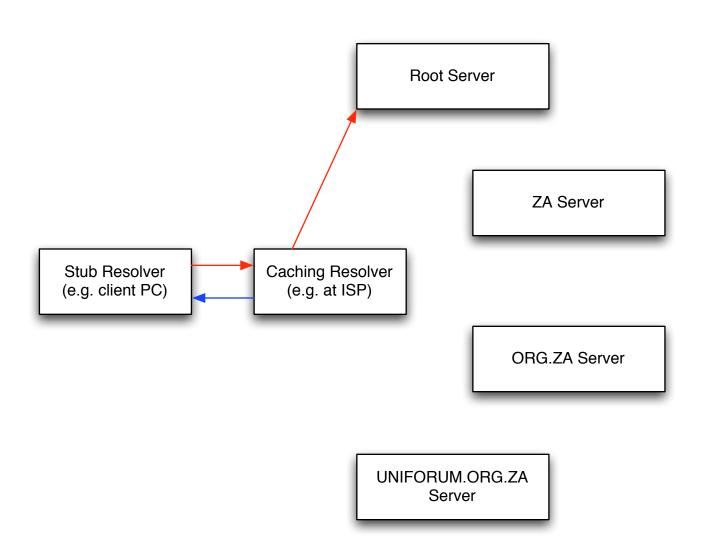
#### Global DNS Failure

- Probability of the entire DNS system failing is low
  - the most important data in the DNS (records which are frequently queried) are cached, usually with high(ish) TTLs
  - the individual root servers are run independently and are under substantial scrutiny
  - coordinated attacks on the root servers tend to be investigated vigorously

## Regional DNS Failure

- If a region becomes partitioned from the Internet, or suffers a prolonged lack of access to the root nameservers for some other reason, the DNS may fail within that region
- Issues affecting small regions do not attract the same attention as issues affecting the whole network
- Regional DNS failure is much more likely than global failure

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#### Loss of Network

- Many countries depend on a relatively non-diverse set of external networks to reach the rest of the world
  - one under-sea cable; one satellite operator
  - a common circuit termination point in a telco hotel somewhere
  - an international network that is close to capacity, and which becomes useless if flooded with junk traffic

# The Distributed F Root Nameserver

#### f.root-servers.net

- Has a single IPv4 address (192.5.5.241)
- Has a single IPv6 address (2001:500::1035)
- Requests sent to those addresses are routed to different nameservers, depending on where the request is made from
  - this behaviour is transparent to devices which send requests to F

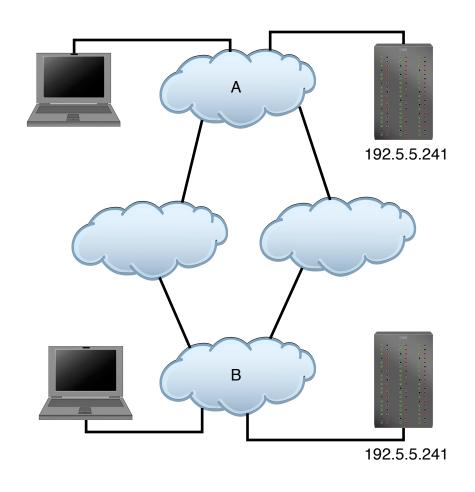
## Unicast, Multicast

- Most traffic on the Internet is unicast
  - packets have a single destination
- Some traffic is multicast
  - packets are directed to multiple destinations

## Anycast

- Traffic to f.root-servers.net is anycast
  - packets are directed to a single instance of F, but different queries (from different places) may land on different instances
  - anycast is identical to unicast from the perspective of the client sending a request

# Anycast Routing



# Hierarchical Anycast

- Some of the F root nameserver nodes provide service to the entire Internet (global nodes)
  - very large, well-connected, secure and over-engineered nodes
- Others provide service to a particular region (local nodes)
  - smaller

# Hierarchical Anycast

- Each local node's routing is organised such that it should not, under normal circumstances, provide service for clients elsewhere in the world
- For more details, see:
  - http://www.isc.org/tn/isc-tn-2003-1.html

## Failure Modes

- If a local node fails, queries to F are automatically routed to a global node
- If a global node fails, queries are automatically routed to another global node
- Catastrophic failure of all global nodes results in continued service by local nodes within their catchment areas

## Failure Modes

- If a region loses international connectivity (e.g. an under-sea cable cut), access to the root nameserver is preserved by virtue of the region's local node
- since the root is reachable, other local nameservers are also reachable (e.g. ZA servers, ORG.ZA servers)
- since TLD servers are reachable, in-country traffic to locally-named services can proceed

### Failure Modes

- A denial of service attack against F launched from outside the region is invisible to users within that region
- A denial of service attack against F launched from within the region is invisible to everybody else in the world
- A widely distributed denial of service attack will cause discomfort proportionate to the size of the region (probably, maybe)

# Triangulation

- Many denial-of-service attacks use sourcespoofed attack traffic
  - time consuming to track back through a network
  - attacks frequently stop before the trace completes
- Watching the relative reactions of local nodes to an attack can help identify the real source

# Logistics and Administrivia

## Sponsorship

- ISC is a non-profit company
- Equipment, colo, networks for remote nodes are paid for by a sponsor
- All equipment is operated exclusively by ISC engineers
- The sponsor covers the ISC's operational costs of running the remote node

# Deployment Status

## Global Nodes

- Palo Alto
- San Francisco

#### Local Nodes

- Amsterdam, Barcelona, Lisbon, Madrid, Moscow, Munich, Paris, Prague, Rome
- São Paulo
- Los Angeles, Monterrey, New York, Ottawa,
   San Jose, Toronto
- Beijing, Dubai, Hong Kong, Jakarta, Osaka, Seoul, Singapore, Taipei, Tel Aviv
- Auckland, Brisbane
- Johannesburg

#### Local Nodes

- Amsterdam, Barcelona, Lisbon, Madrid, Moscow, Munich, Paris, Prague, Rome
- São Paulo
- Los Angeles, Monterrey, New York, Ottawa,
   San Jose, Toronto
- Beijing, Dubai, Hong Kong, Jakarta, Osaka, Seoul, Singapore, Taipei, Tel Aviv
- Auckland, Brisbane
- Johannesburg, **Nairobi**

## The Nairobi F

#### Vital Statistics

- Physically colocated with the KIXP switch
- I00 Mbit/s connection to the KIXP
- Two redundant, much lower-capacity transit paths via two independent ISPs for management, measurement, zone transfers
- Cluster of two nameservers sharing the query load

## Using the Local F

- You may be already using it
  - traceroute f.root-servers.net
  - dig @f.root-servers.net hostname.bind chaos txt

## Before...

```
[halibut:~]$ traceroute f.root-servers.net
traceroute to f.root-servers.net (192.5.5.241), 64 hops max, 40 byte packets
   router.cctld.or.ke (196.216.0.62) 1.945 ms 7.147 ms 1.165 ms
   196.216.66.5 (196.216.66.5) 44.967 ms 23.918 ms 12.420 ms
   217.21.112.4.swiftkenya.com (217.21.112.4) 5.141 ms 9.491 ms 5.791 ms
   193.220.225.5 (193.220.225.5) 8.919 ms 5.708 ms 5.898 ms
   no-nit-tn-7.taide.net (193.219.192.7) 538.820 ms 539.738 ms 550.056 ms
   no-nit-tn-5.taide.net (193.219.193.145) 540.073 ms 551.002 ms 536.818 ms
   pos5-1.qw3.osl2.alter.net (146.188.39.1) 535.738 ms 536.197 ms 534.790 ms
   so-3-0-0.xr2.osl2.alter.net (146.188.15.97) 535.701 ms 542.140 ms 543.969 ms
   so-4-2-0.tr1.stk2.alter.net (146.188.15.61) 541.221 ms 545.562 ms 544.435 ms
   so-7-0-0.ir2.dca4.alter.net (146.188.11.226) 653.929 ms 652.082 ms 649.199 ms
10
11
   so-1-0-0.il2.dca6.alter.net (146.188.13.45) 658.517 ms 652.177 ms 664.978 ms
   0.so-0-2-0.tl2.sac1.alter.net (152.63.0.190) 887.784 ms 739.093 ms 717.126 ms
12
   0.so-1-3-0.xl2.pao1.alter.net (152.63.48.181) 718.044 ms 720.835 ms 727.418 ms
13
   pos1-0.xr2.pao1.alter.net (152.63.54.78) 717.283 ms 716.201 ms 714.212 ms
14
   188.atm7-0.gw10.pao1.alter.net (152.63.53.21) 778.208 ms 731.906 ms 832.482 ms
15
   isc-pao-gw.customer.alter.net (157.130.205.230) 717.801 ms 712.912 ms 712.718 ms
16
   f.root-servers.net (192.5.5.241) 743.804 ms 721.633 ms 746.818 ms
Fhalibut:~1$
```

#### ... and After

```
[halibut:~]$ traceroute f.root-servers.net
traceroute to f.root-servers.net (199.6.6.14), 64 hops max, 40 byte packets
1 router.cctld.or.ke (196.216.0.62) 244.241 ms 1.159 ms 1.099 ms
2 196.216.66.5 (196.216.66.5) 8.678 ms 4.942 ms 31.862 ms
3 80.240.202.54.swiftkenya.com (80.240.202.54) 22.455 ms 15.803 ms 14.864 ms
4 198.32.143.125 (198.32.143.125) 40.770 ms 7.192 ms 7.786 ms
5 f.root-servers.net (192.5.5.241) 10.906 ms 10.894 ms *
[halibut:~]$
```

# Sponors

• KENIC