

IP and Networking Basics



Ubuntunet Workshop 2012

Slides from SIE@AFNOG

Internet History

1961-1972: Early packet-switching principles

- ❑ 1961: Kleinrock - queueing theory shows effectiveness of packet-switching
- ❑ 1964: Baran - packet-switching in military nets
- ❑ 1967: ARPAnet conceived by Advanced Research Projects Agency
- ❑ 1969: first ARPAnet node operational

1972:

ARPAnet demonstrated publicly
NCP (Network Control Protocol) first host-host protocol
first e-mail program
ARPAnet has 15 nodes

Internet History

1972-1980: Internetworking, new and proprietary nets

- ❑ 1970: ALOHAnet satellite network in Hawaii
- ❑ 1973: Metcalfe's PhD thesis proposes Ethernet
- ❑ 1974: Cerf and Kahn - architecture for interconnecting networks
- ❑ Late 70's: proprietary architectures: DECnet, SNA, XNA
- ❑ late 70's: switching fixed length packets (ATM precursor)
- ❑ 1979: ARPAnet has 200 nodes

Cerf and Kahn's
internetworking principles:
minimalism, autonomy - no
internal changes required
to interconnect networks
best effort service model
stateless routers
decentralized control
define today's Internet
architecture

Internet History

1980-1990: new protocols, a proliferation of networks

- ❑ 1983: deployment of TCP/IP
- ❑ 1982: SMTP e-mail protocol defined
- ❑ 1983: DNS defined for name-to-IP-address translation
- ❑ 1985: FTP protocol defined
- ❑ 1988: TCP congestion control

New national networks:
Csnet, BITnet, NSFnet,
Minitel

100,000 hosts connected to
confederation of networks

Internet History

1990, 2000's: commercialisation, the Web, new apps

- ❑ Early 1990's: ARPAnet decommissioned
- ❑ 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- ❑ early 1990s: Web
 - hypertext [Bush 1945, Nelson 1960's]
 - HTML, HTTP: Berners-Lee
 - 1994: Mosaic, later Netscape
 - late 1990's: commercialization of the Web

Late 1990's – 2000's:

more killer apps: instant messaging, peer2peer file sharing (e.g., Napster)
network security to forefront
est. 50 million host, 100 million+ users
backbone links running at Gbps

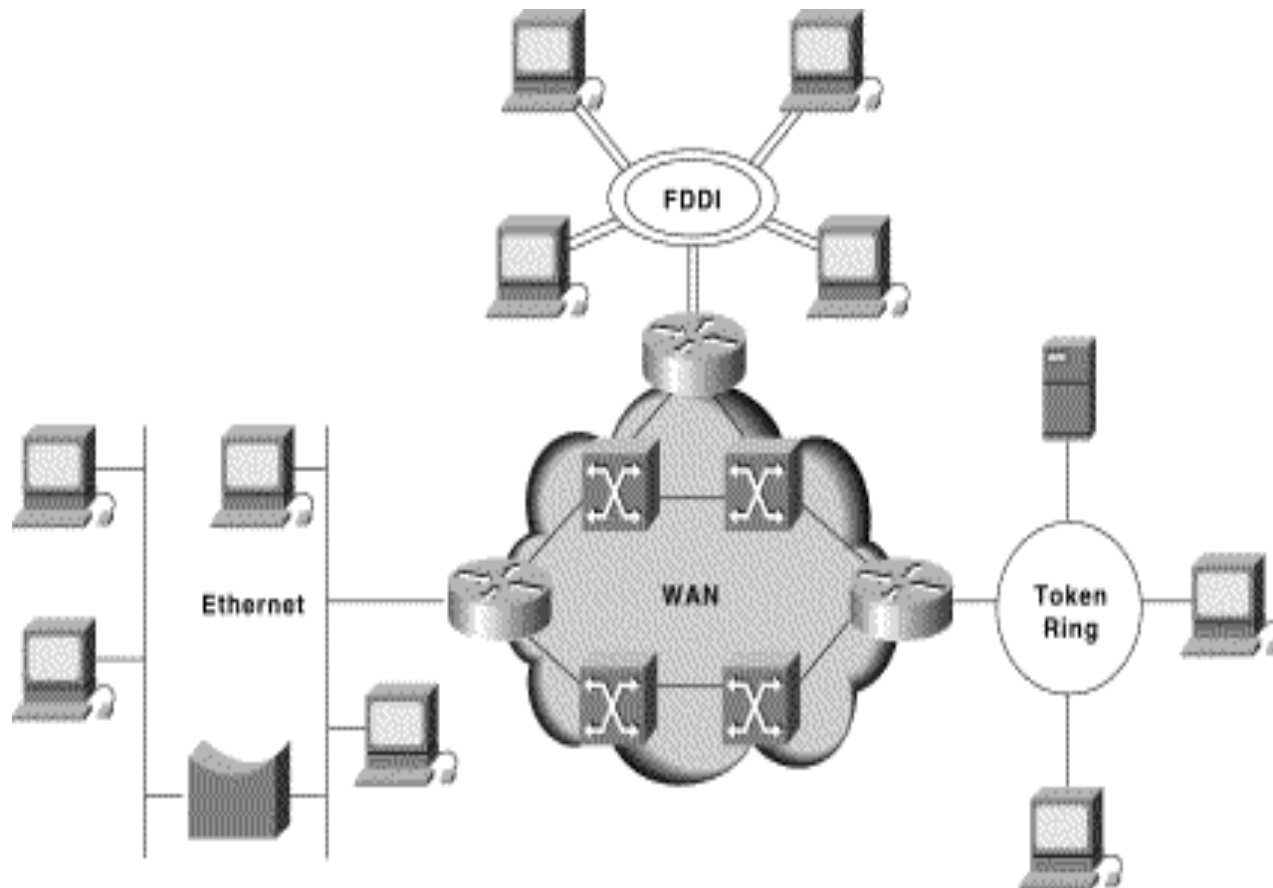
now: 40-100 Gbps

youtube, social networking
depletion of Ipv4 address space

The (capital “I”) Internet

- ❑ The world-wide network of TCP/IP networks
- ❑ Different people or organisations own different parts
- ❑ Different parts use different technologies
- ❑ Interconnections between the parts
- ❑ Interconnections require agreements
 - sale/purchase of service
 - contracts
 - “peering” agreements
- ❑ No central control or management

A small internetwork or (small “i”) “internet”

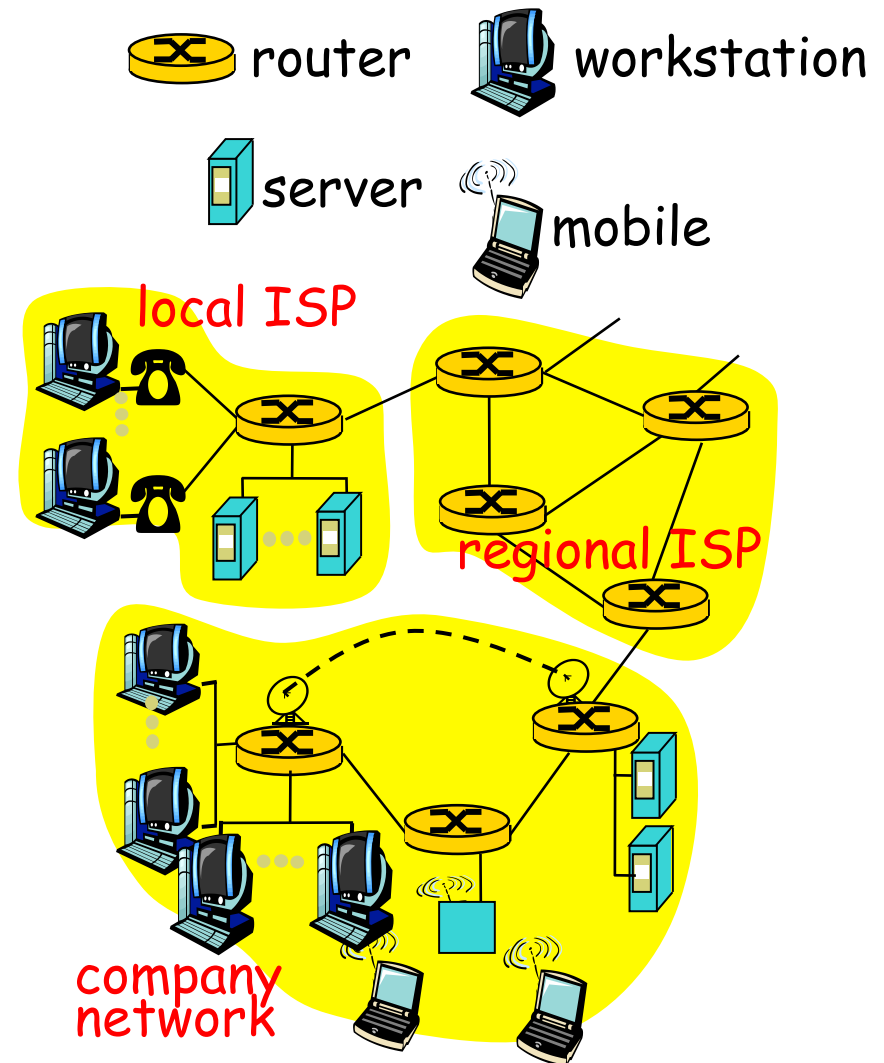


The principle of “Internetworking”

- ❑ We have lots of little networks
- ❑ Many different owners/operators
- ❑ Many different types
 - Ethernet, dedicated leased lines, dialup, optical, broadband, wireless, ...
- ❑ Each type has its own idea of low level addressing and protocols
- ❑ We want to connect them all together and provide a unified view of the whole lot (treat the collection of networks as a single large internetwork)

What is the Internet: “nuts and bolts” view

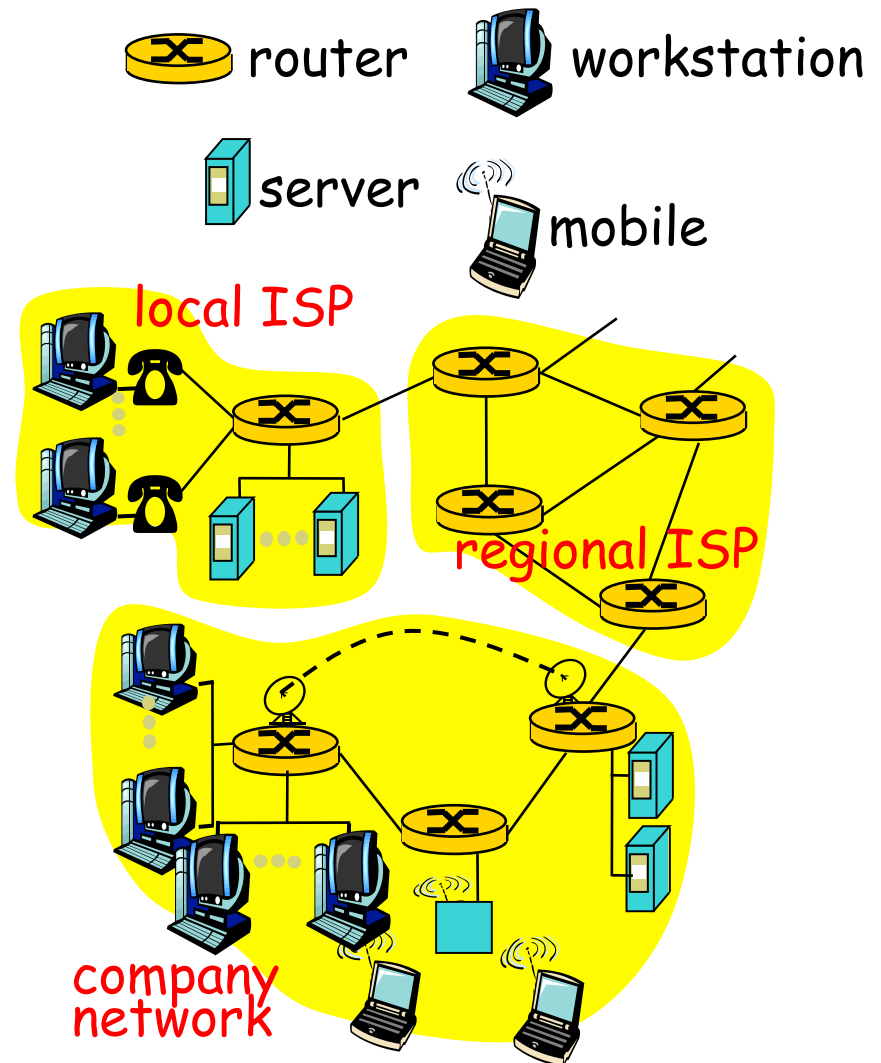
- ❑ millions of connected computing devices: hosts, end-systems
 - PC's workstations, servers
 - PDA's phones, toasters
 - running network apps
- ❑ communication links
 - fiber, copper, radio, satellite
- ❑ routers: forward packets (chunks) of data through network



What is the Internet:

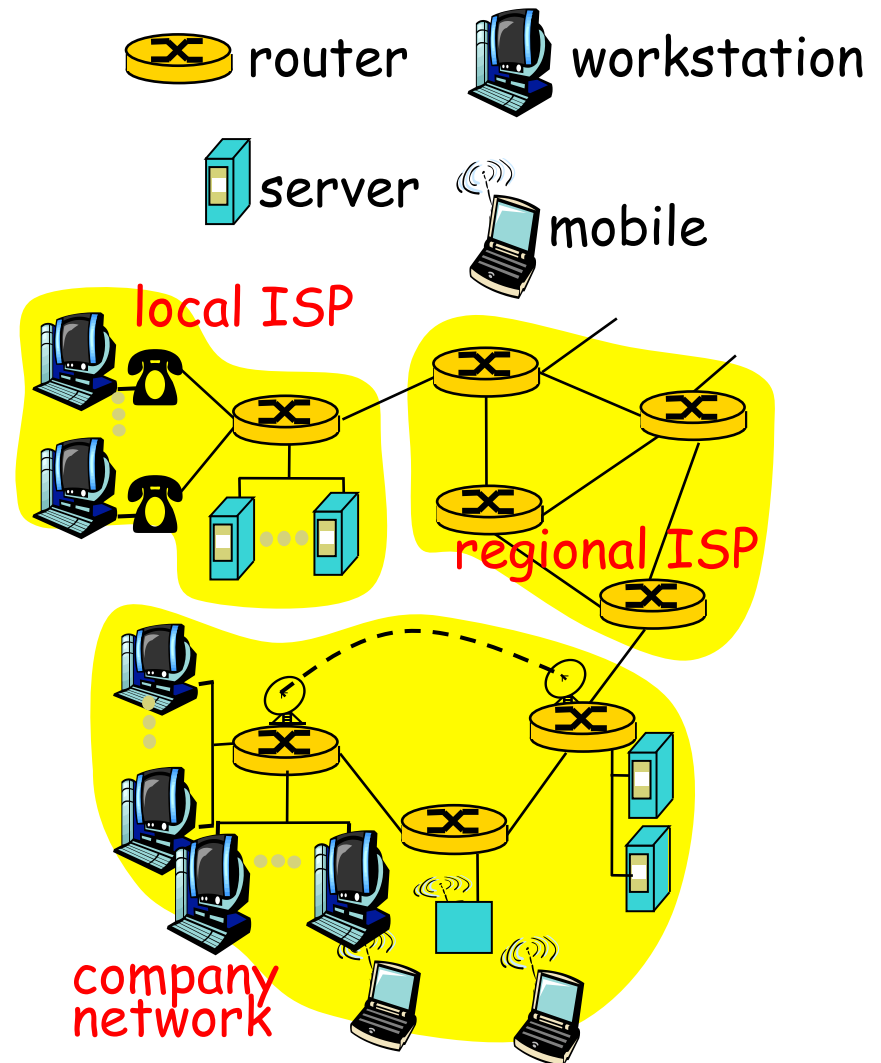
“nuts and bolts” view

- protocols: control sending, receiving of messages
 - e.g., TCP, IP, HTTP, FTP, PPP
- Internet: “network of networks”
 - loosely hierarchical
 - public Internet versus private intranet
- Internet standards
 - RFC: Request for comments
 - IETF: Internet Engineering Task Force



What is the Internet: a service view

- communication infrastructure enables distributed applications:
 - WWW, email, games, e-commerce, database, e-voting, more?
- communication services provided:
 - connectionless
 - connection-oriented



Connectionless Paradigm

- ❑ There is no “connection” in IP
 - Packets can be delivered out-of-order
 - Each packet can take a different path to the destination
 - No error detection or correction in payload
 - No congestion control (beyond “drop”)
- ❑ TCP mitigates these for connection-oriented applications
 - error correction is by retransmission
 - Packet drops as congestion signalling

OSI Stack & TCP/IP Architecture

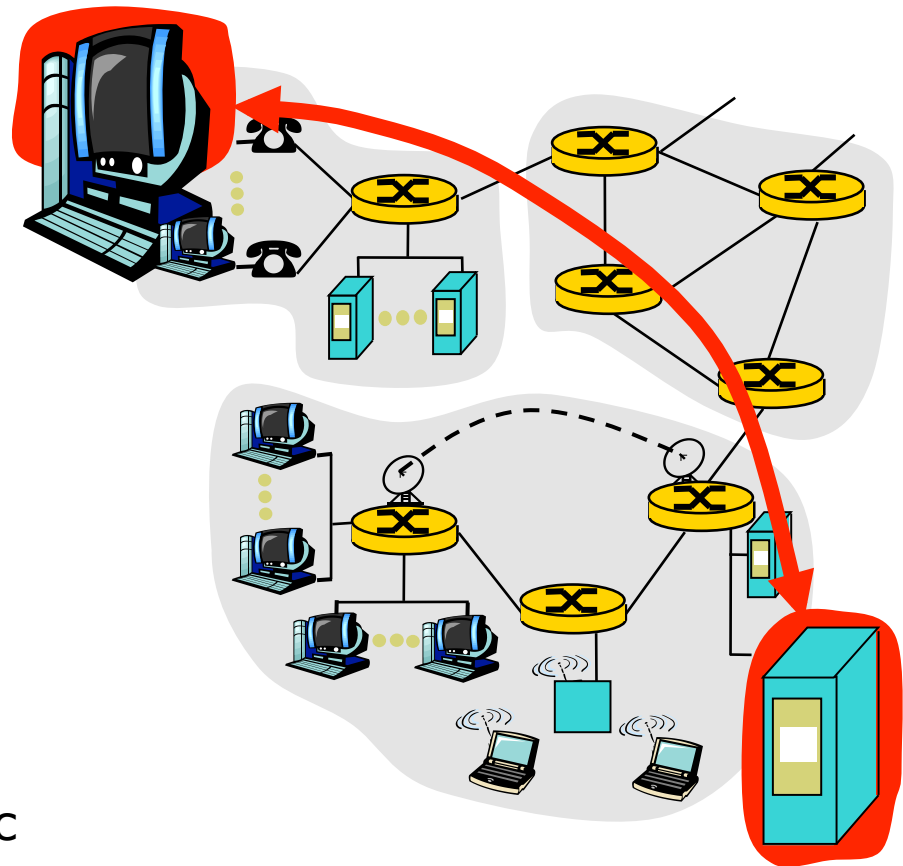


Principles of the Internet

- ❑ Edge vs. core (end-systems vs. routers)
 - Dumb network
 - Intelligence at the end-systems
- ❑ Different communication paradigms
 - Connection oriented vs. connection less
 - Packet vs. circuit switching
- ❑ Layered System
- ❑ Network of collaborating networks

The network edge

- ❑ end systems (hosts):
 - run application programs
 - e.g., WWW, email
 - at “edge of network”
- ❑ client/server model:
 - client host requests, receives service from server
 - e.g., WWW client (browser)/server; email client/server
- ❑ peer-peer model:
 - host interaction symmetric e.g.: teleconferencing



Network edge: connection-oriented service

- Goal: data transfer between end sys.
- handshaking: setup (prepare for) data transfer ahead of time
 - Hello, hello back human protocol
 - set up “state” in two communicating hosts
- TCP - Transmission Control Protocol
 - Internet’s connection-oriented service

TCP service [RFC 793]

reliable, in-order byte-stream data transfer

loss: acknowledgements and retransmissions

flow control:

sender won’t overwhelm receiver

congestion control:

senders “slow down sending rate” when network congested

Network edge: connectionless service

- Goal: data transfer between end systems
- UDP - User Datagram Protocol [RFC 768]: Internet's connectionless service
 - unreliable data transfer
 - no flow control
 - no congestion control

Protocol “Layers”

- ❑ Networks are complex!
- ❑ many “pieces”:
 - hosts
 - routers
 - links of various media
 - applications
 - protocols
 - hardware, software

Question:

Is there any hope of
organizing structure of
network?

Or at least in our
discussion of
networks?

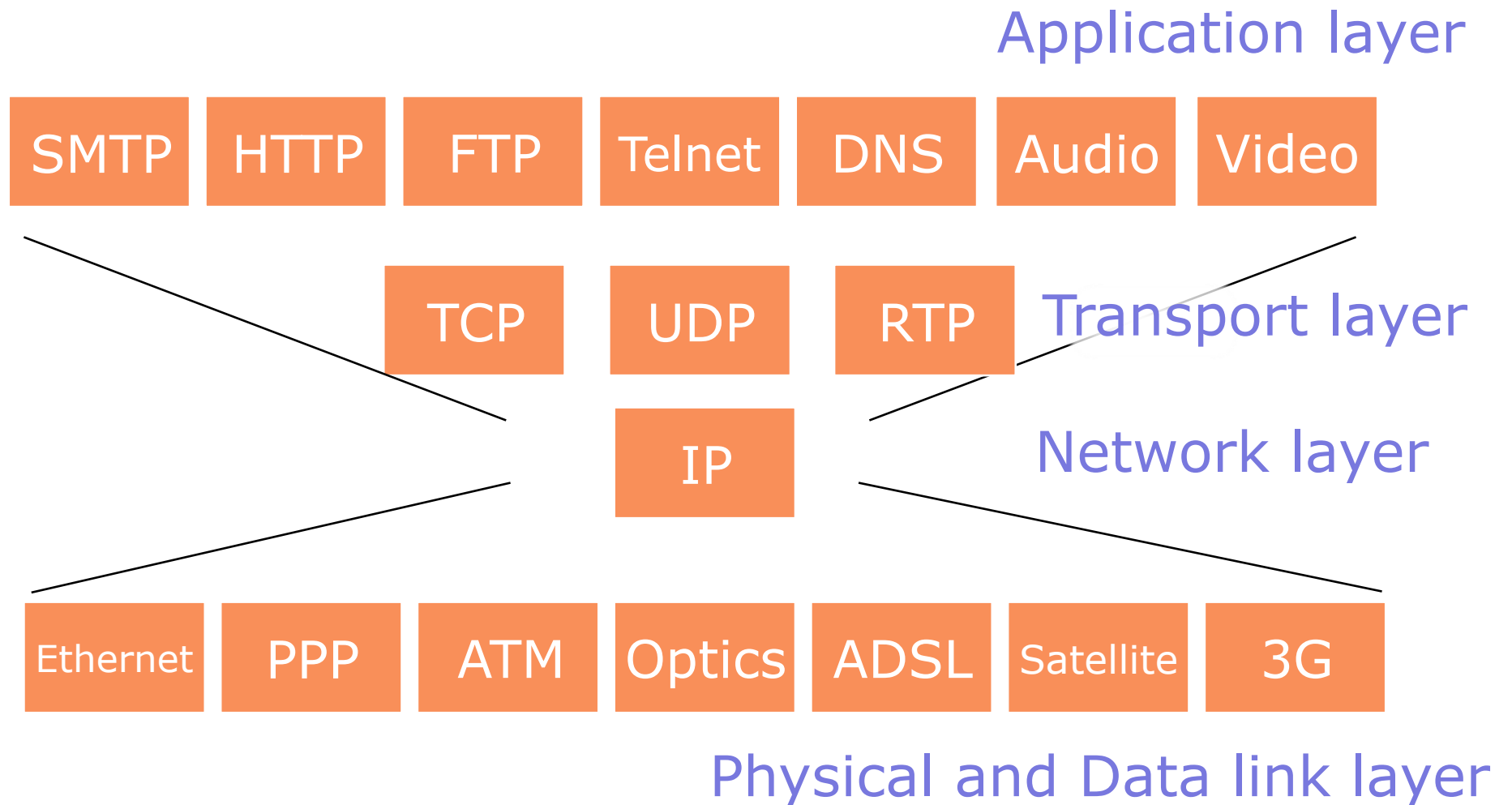
The unifying effect of the network layer

- ❑ Define a protocol that works in the same way with any underlying network
- ❑ Call it the network layer (e.g. IP)
- ❑ IP routers operate at the network layer
- ❑ IP over anything
- ❑ Anything over IP

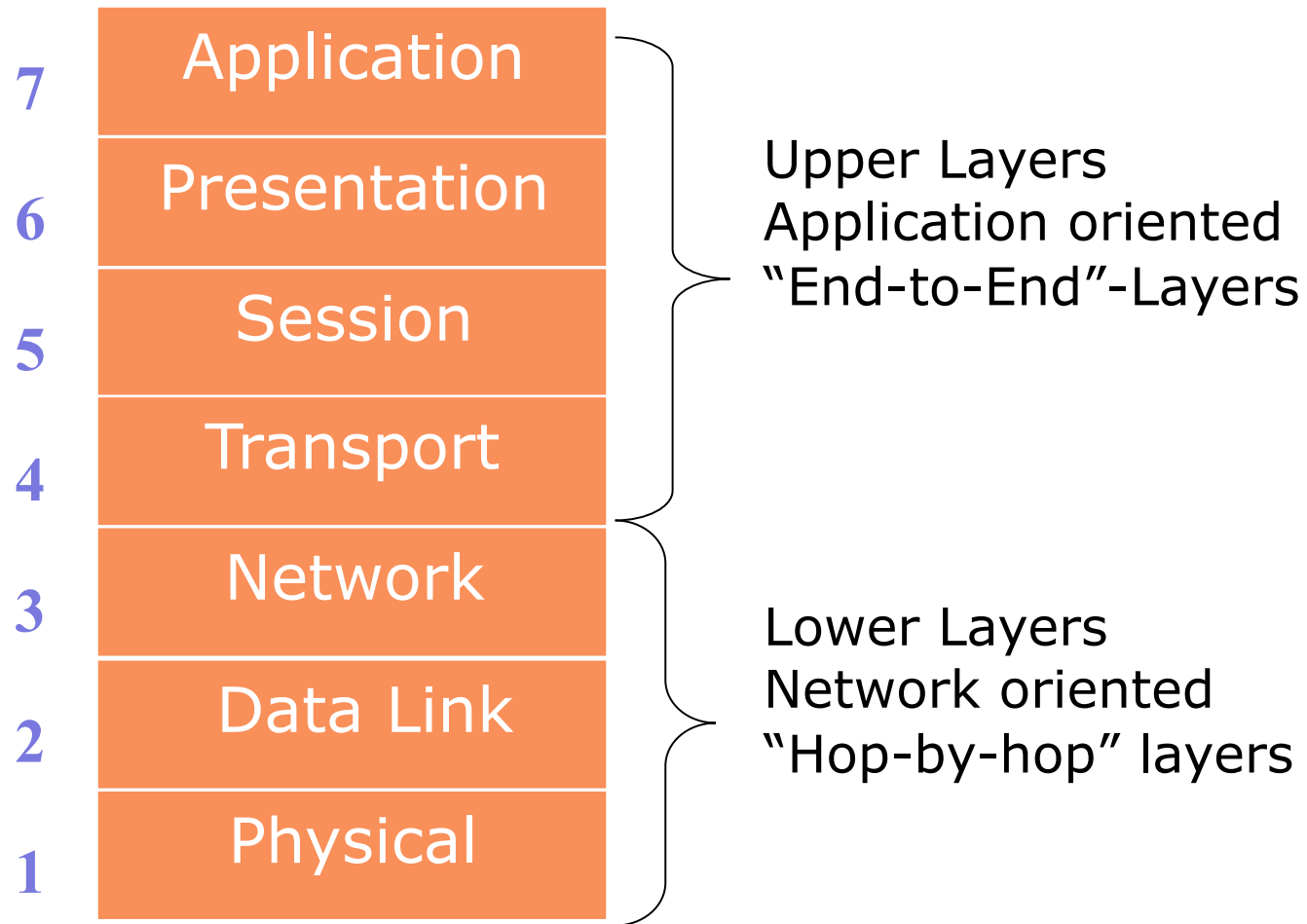
Why layering?

- ❑ Dealing with complex systems:
- ❑ explicit structure allows identification, relationship of complex system's pieces
 - layered reference model for discussion
- ❑ Modularisation eases maintenance, updating of system
 - change of implementation of layer's service transparent to rest of system
 - e.g., change in gate procedure does not affect rest of system

The IP Hourglass Model



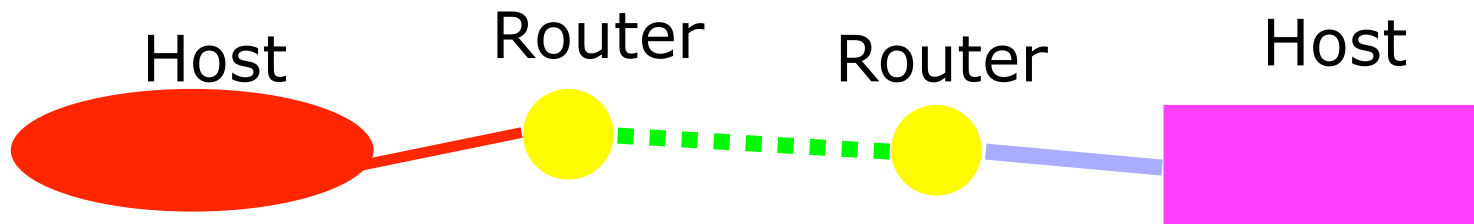
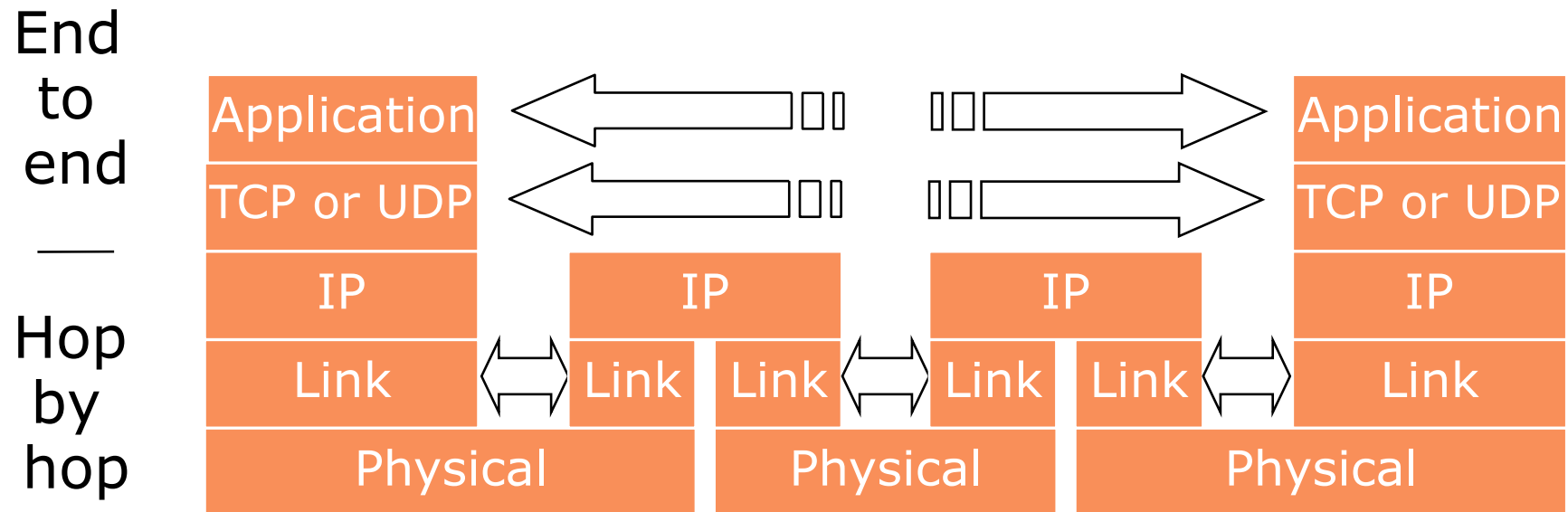
The OSI Model



OSI Model and the Internet

- ❑ Internet protocols are not directly based on the OSI model
- ❑ However, we do often use the OSI numbering system. You should at least remember these:
 - Layer 7: Application
 - Layer 4: Transport (e.g. TCP, UDP)
 - Layer 3: Network (IP)
 - Layer 2: Data link
 - Layer 1: Physical

Layer Interaction: TCP/IP Model



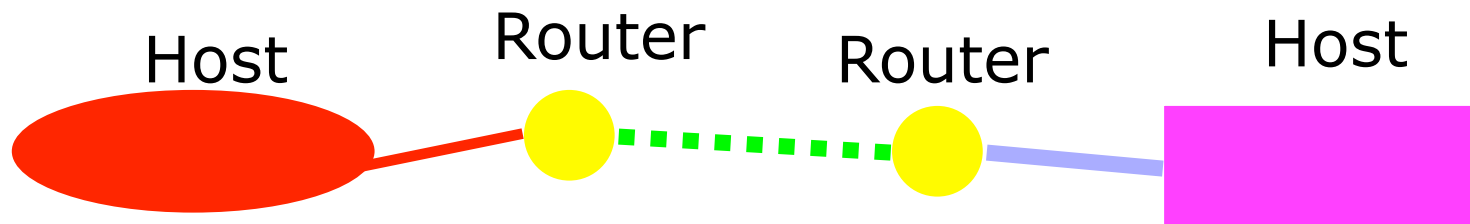
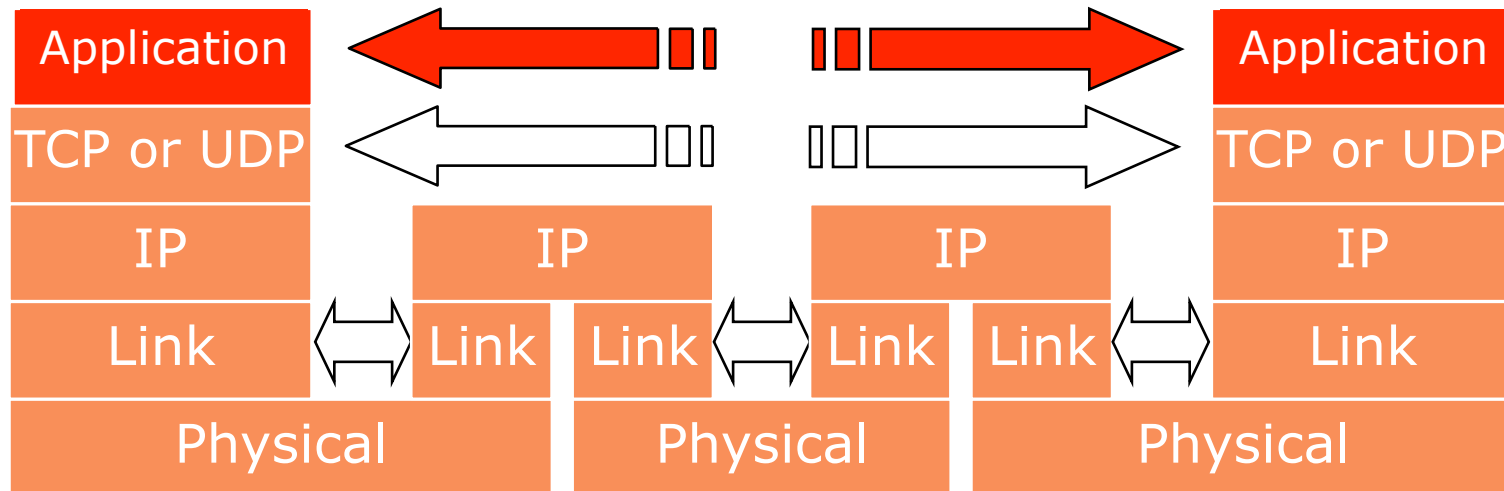
End-to-end layers

- Upper layers are “end-to-end”
- Applications at the two ends behave as if they can talk directly to each other
- They do not concern themselves with the details of what happens in between

Hop-by-hop layers

- ❑ At the lower layers, devices share access to the same physical medium
- ❑ Devices communicate directly with each other
- ❑ The network layer (IP) has some knowledge of how many small networks are interconnected to make a large internet
- ❑ Information moves one hop at a time, getting closer to the destination at each hop

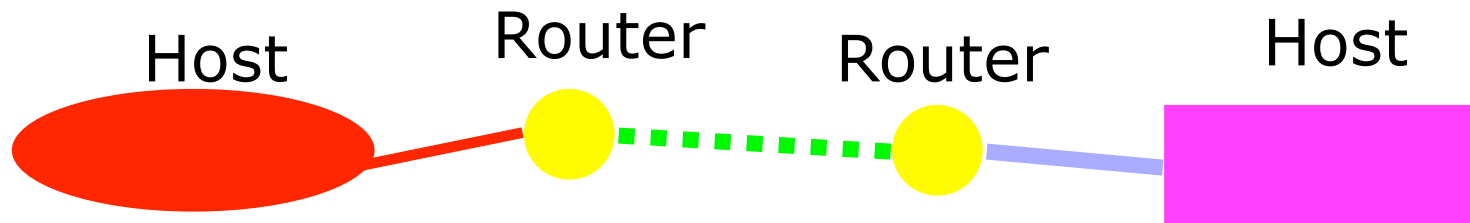
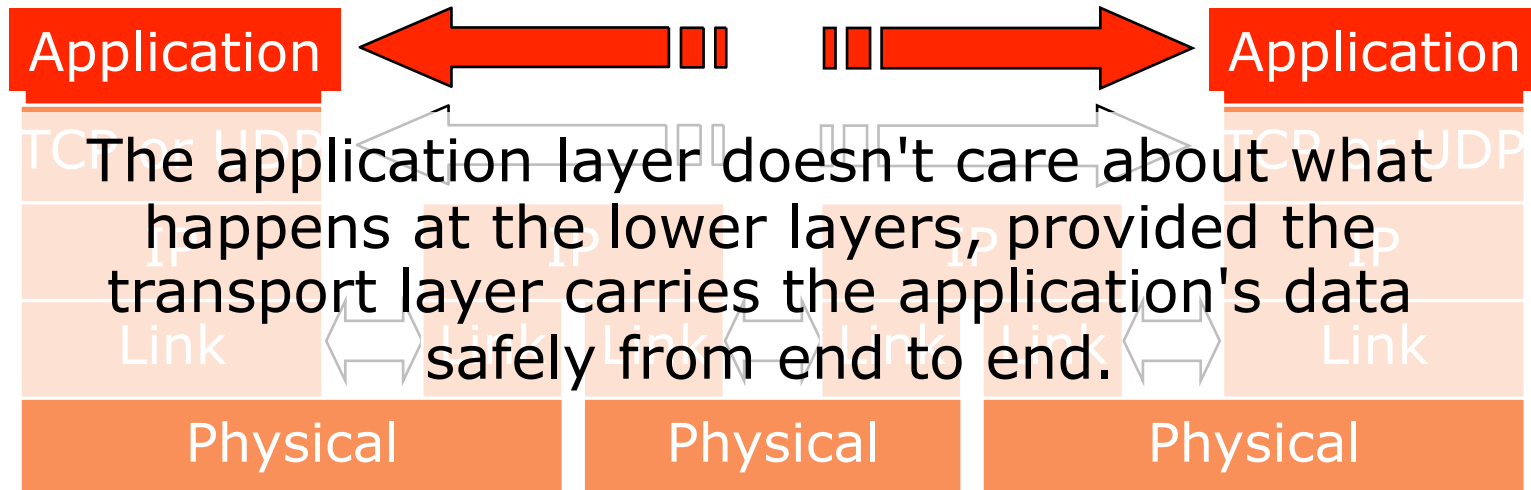
Layer Interaction: TCP/IP Model



Layer Interaction:

The Application Layer

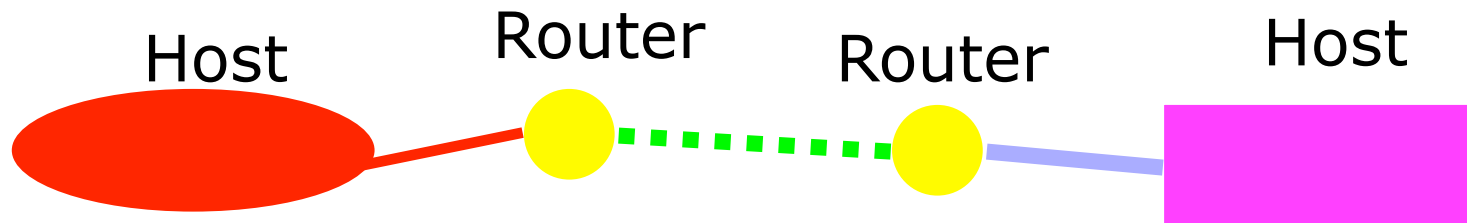
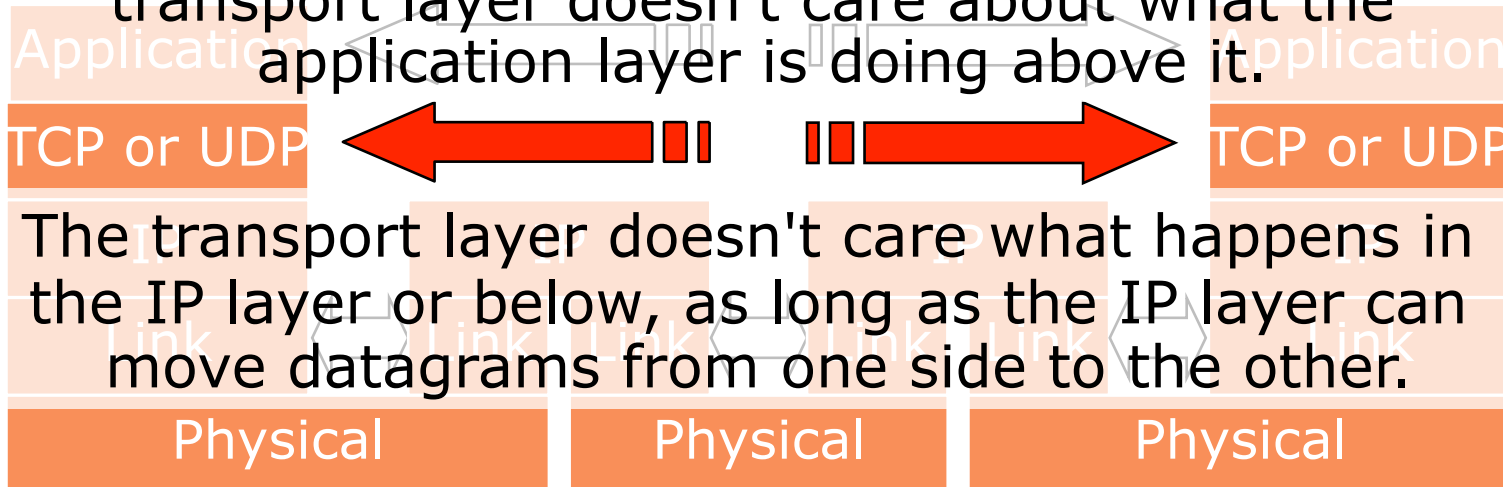
Applications behave as if they can talk to each other, but in reality the application at each side talks to the TCP or UDP service below it.



Layer Interaction:

The Transport Layer

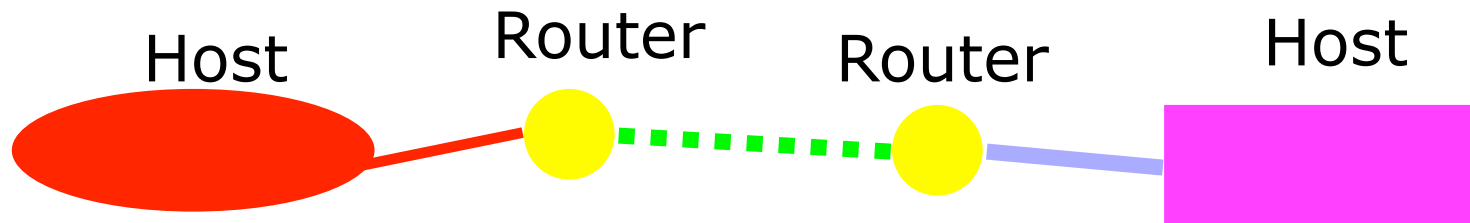
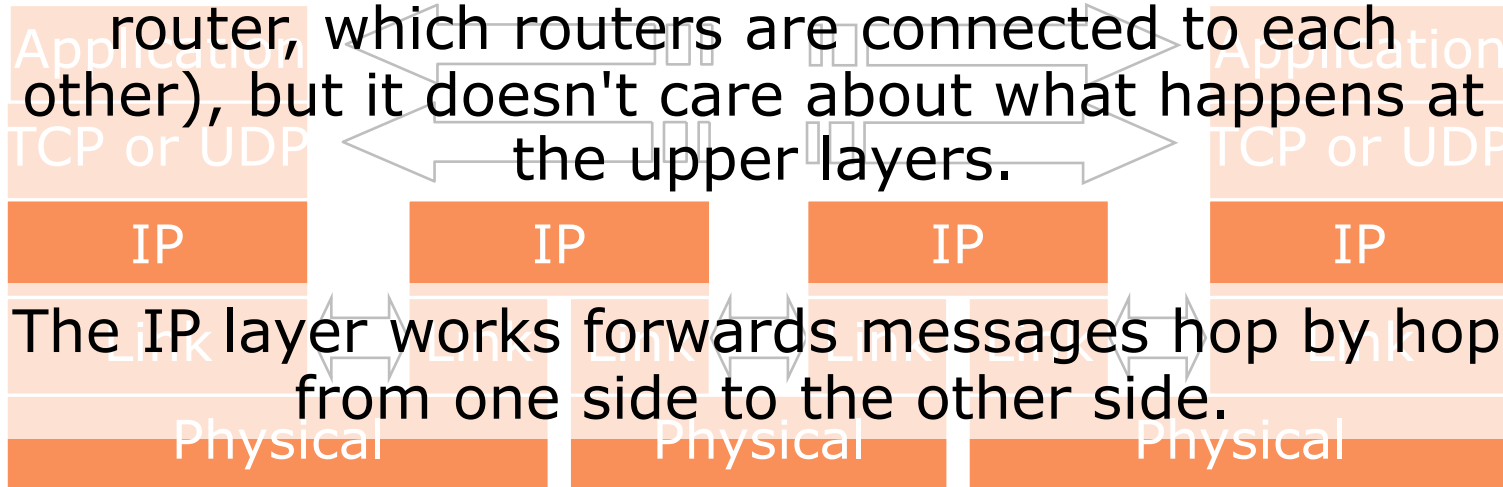
The transport layer instances at the two ends act as if they are talking to each other, but in reality they are each talking to the IP layer below it. The transport layer doesn't care about what the application layer is doing above it.



Layer Interaction:

The Network Layer (IP)

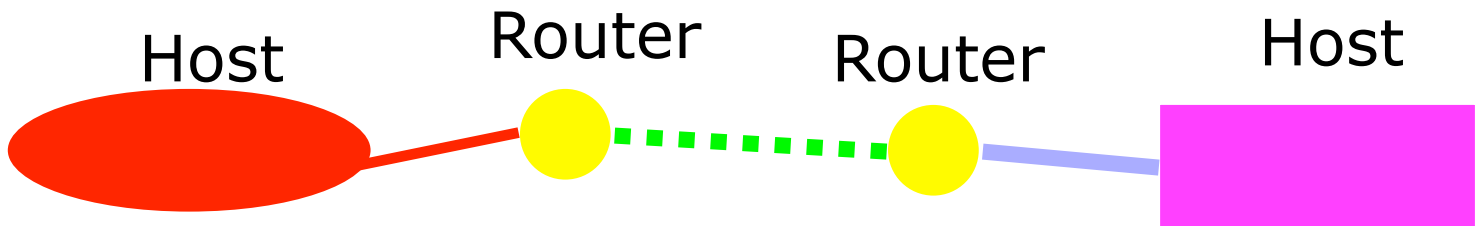
The IP layer has to know a lot about the topology of the network (which host is connected to which router, which routers are connected to each other), but it doesn't care about what happens at the upper layers.



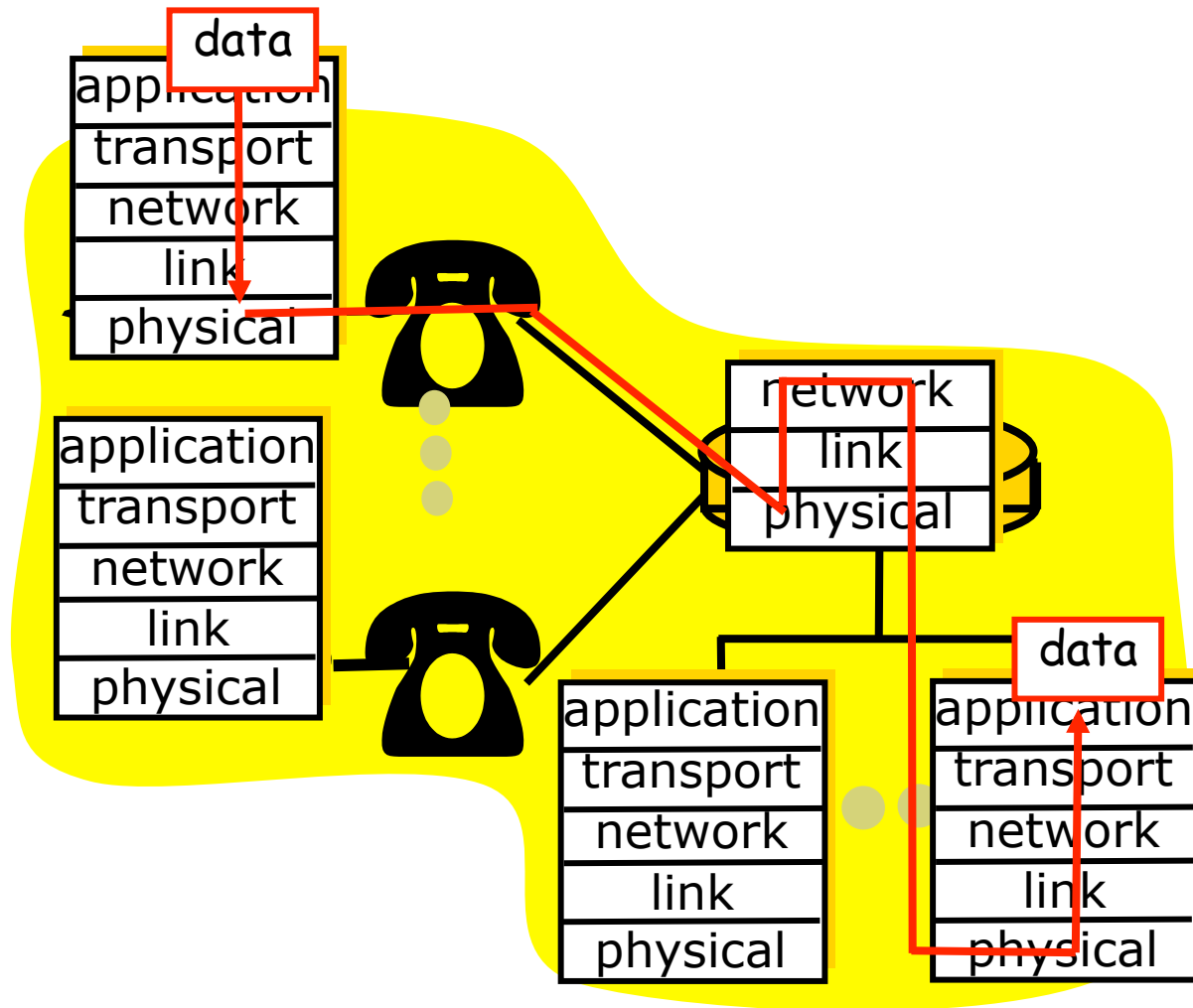
Age Group	Percentage
18-24	15%
25-34	25%
35-44	30%
45-54	20%
55-64	10%
65-74	5%
75-84	10%
85+	5%

but it is very closely tied to the physical layer below it.

All links are independent of each other, and have no way of communicating with each other.



Layering: physical communication

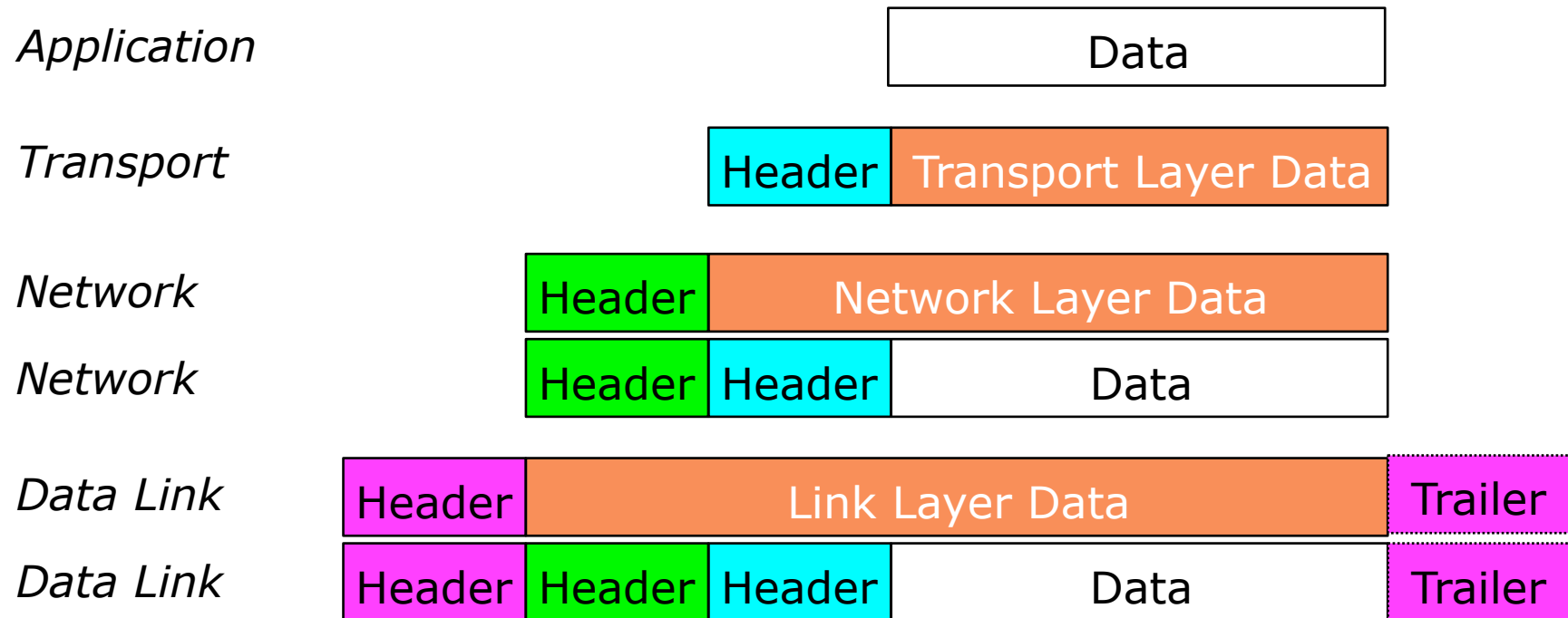


Frame, Datagram, Segment, Packet

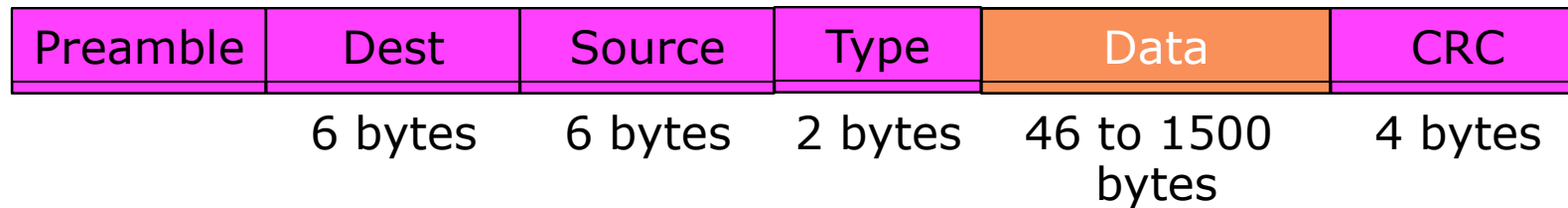
- Different names for packets at different layers
 - Ethernet (link layer) frame
 - IP (network layer) datagram
 - TCP (transport layer) segment
- Terminology is not strictly followed
 - we often just use the term “packet” at any layer

Encapsulation & Decapsulation

- Lower layers add headers (and sometimes trailers) to data from higher layers



Layer 2 - Ethernet frame



- ❑ Destination and source are 48-bit MAC addresses (e.g., 00:26:4a:18:f6:aa)
- ❑ Type 0x0800 means that the “data” portion of the Ethernet frame contains an IPv4 datagram. Type 0x0806 for ARP. Type 0x86DD for IPv6.
- ❑ “Data” part of layer 2 frame contains a layer 3 datagram.

Layer 3 - IPv4 datagram

Version	IHL	Diff Services	Total Length	
Identification			Flags	Fragment Offset
Time to Live		Protocol	Header Checksum	
Source Address (32-bit IPv4 address)				
Destination Address (32-bit IPv4 address)				
Options				Padding
Data (contains layer 4 segment)				

Version = 4

If no options, IHL = 5
Source and
Destination are 32-bit
IPv4 addresses

▣ Protocol = 6 means
data portion contains
a TCP segment.
Protocol = 17 means
UDP.

Layer 4 - TCP segment

Source Port							Destination Port						
Sequence Number													
Acknowledgement Number													
Data Offset	Reserved	U	A	E	R	S	F	Window					
		R	C	O	S	Y	I						
		G	K	L	T	N	N						
Checksum							Urgent Pointer						
Options										Padding			
Data (contains application data)													

- Source and Destination are 16-bit TCP port numbers (IP addresses are implied by the IP header)
- If no options, Data Offset = 5 (which means 20 octets)



Questions?
